GNU libiberty

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1 Using

To date, `libiberty` is generally not installed on its own. It has evolved over years but does not have its own version number nor release schedule.

Possibly the easiest way to use `libiberty` in your projects is to drop the `libiberty` code into your project's sources, and to build the library along with your own sources; the library would then be linked in at the end. This prevents any possible version mismatches with other copies of libiberty elsewhere on the system.

Passing `--enable-install-libiberty` to the `configure` script when building `libiberty` causes the header files and archive library to be installed when *`make install`* is run. This option also takes an (optional) argument to specify the installation location, in the same manner as `--prefix`.

For your own projects, an approach which offers stability and flexibility is to include `libiberty` with your code, but allow the end user to optionally choose to use a previously-installed version instead. In this way the user may choose (for example) to install `libiberty` as part of GCC, and use that version for all software built with that compiler. (This approach has proven useful with software using the GNU `readline` library.)

Making use of `libiberty` code usually requires that you include one or more header files from the `libiberty` distribution. (They will be named as necessary in the function descriptions.) At link time, you will need to add `-liberty` to your link command invocation.

2 Overview

Functions contained in `libiberty` can be divided into three general categories.

2.1 Supplemental Functions

Certain operating systems do not provide functions which have since become standardized, or at least common. For example, the Single Unix Specification Version 2 requires that the `basename` function be provided, but an OS which predates that specification might not have this function. This should not prevent well-written code from running on such a system.

Similarly, some functions exist only among a particular "flavor" or "family" of operating systems. As an example, the `bzero` function is often not present on systems outside the BSD-derived family of systems.

Many such functions are provided in `libiberty`. They are quickly listed here with little description, as systems which lack them become less and less common. Each function *foo* is implemented in `foo.c` but not declared in any `libiberty` header file; more comments and caveats for each function's implementation are often available in the source file. Generally, the function can simply be declared as `extern`.

2.2 Replacement Functions

Some functions have extremely limited implementations on different platforms. Other functions are tedious to use correctly; for example, proper use of `malloc` calls for the return value to be checked and appropriate action taken if memory has been exhausted. A group of "replacement functions" is available in `libiberty` to address these issues for some of the most commonly used subroutines.

All of these functions are declared in the `libiberty.h` header file. Many of the implementations will use preprocessor macros set by GNU Autoconf, if you decide to make use of that program. Some of these functions may call one another.

2.2.1 Memory Allocation

The functions beginning with the letter '`x`' are wrappers around standard functions; the functions provided by the system environment are called and their results checked before the results are passed back to client code. If the standard functions fail, these wrappers will terminate the program. Thus, these versions can be used with impunity.

2.2.2 Exit Handlers

The existence and implementation of the `atexit` routine varies amongst the flavors of Unix. `libiberty` provides an unvarying dependable implementation via `xatexit` and `xexit`.

2.2.3 Error Reporting

These are a set of routines to facilitate programming with the system `errno` interface. The `libiberty` source file `strerror.c` contains a good deal of documentation for these functions.

2.3 Extensions

`libiberty` includes additional functionality above and beyond standard functions, which has proven generically useful in GNU programs, such as obstacks and regex. These functions are often copied from other projects as they gain popularity, and are included here to provide a central location from which to use, maintain, and distribute them.

2.3.1 Obstacks

An *obstack* is a pool of memory containing a stack of objects. You can create any number of separate obstacks, and then allocate objects in specified obstacks. Within each obstack, the last object allocated must always be the first one freed, but distinct obstacks are independent of each other.

Aside from this one constraint of order of freeing, obstacks are totally general: an obstack can contain any number of objects of any size. They are implemented with macros, so allocation is usually very fast as long as the objects are usually small. And the only space overhead per object is the padding needed to start each object on a suitable boundary.

2.3.1.1 Creating Obstacks

The utilities for manipulating obstacks are declared in the header file `obstack.h`.

`struct obstack` [Data Type]

> An obstack is represented by a data structure of type `struct obstack`. This structure has a small fixed size; it records the status of the obstack and how to find the space in which objects are allocated. It does not contain any of the objects themselves. You should not try to access the contents of the structure directly; use only the macros described in this chapter.

You can declare variables of type `struct obstack` and use them as obstacks, or you can allocate obstacks dynamically like any other kind of object. Dynamic allocation of obstacks allows your program to have a variable number of different stacks. (You can even allocate an obstack structure in another obstack, but this is rarely useful.)

All the macros that work with obstacks require you to specify which obstack to use. You do this with a pointer of type `struct obstack *`. In the following, we often say "an obstack" when strictly speaking the object at hand is such a pointer.

The objects in the obstack are packed into large blocks called *chunks*. The `struct obstack` structure points to a chain of the chunks currently in use.

The obstack library obtains a new chunk whenever you allocate an object that won't fit in the previous chunk. Since the obstack library manages chunks automatically, you don't need to pay much attention to them, but you do need to supply a function which the obstack library should use to get a chunk. Usually you supply a function which uses `malloc` directly or indirectly. You must also supply a function to free a chunk. These matters are described in the following section.

2.3.1.2 Preparing for Using Obstacks

Each source file in which you plan to use obstacks must include the header file `obstack.h`, like this:

```
#include <obstack.h>
```

Also, if the source file uses the macro `obstack_init`, it must declare or define two macros that will be called by the obstack library. One, `obstack_chunk_alloc`, is used to allocate the chunks of memory into which objects are packed. The other, `obstack_chunk_free`, is used to return chunks when the objects in them are freed. These macros should appear before any use of obstacks in the source file.

Usually these are defined to use `malloc` via the intermediary `xmalloc` (see Section "Unconstrained Allocation" in *The GNU C Library Reference Manual*). This is done with the following pair of macro definitions:

```
#define obstack_chunk_alloc xmalloc
#define obstack_chunk_free free
```

Though the memory you get using obstacks really comes from `malloc`, using obstacks is faster because `malloc` is called less often, for larger blocks of memory. See Section 2.3.1.10 [Obstack Chunks], page 11, for full details.

At run time, before the program can use a `struct obstack` object as an obstack, it must initialize the obstack by calling `obstack_init` or one of its variants, `obstack_begin`, `obstack_specify_allocation`, or `obstack_specify_allocation_with_arg`.

`int obstack_init (struct obstack *`*`obstack-ptr`*`)` [Function]

Initialize obstack *obstack-ptr* for allocation of objects. This macro calls the obstack's `obstack_chunk_alloc` function. If allocation of memory fails, the function pointed to by `obstack_alloc_failed_handler` is called. The `obstack_init` macro always returns 1 (Compatibility notice: Former versions of obstack returned 0 if allocation failed).

Here are two examples of how to allocate the space for an obstack and initialize it. First, an obstack that is a static variable:

```
static struct obstack myobstack;
...
obstack_init (&myobstack);
```

Second, an obstack that is itself dynamically allocated:

```
struct obstack *myobstack_ptr
  = (struct obstack *) xmalloc (sizeof (struct obstack));

obstack_init (myobstack_ptr);
```

`int obstack_begin (struct obstack *`*`obstack-ptr`*`, size_t chunk_size)` [Function]

Like `obstack_init`, but specify chunks to be at least *chunk_size* bytes in size.

`int obstack_specify_allocation (struct obstack *`*`obstack-ptr`*`, size_t chunk_size, size_t alignment, void *(*chunkfun) (size_t), void (*freefun) (void *))` [Function]

Like `obstack_init`, specifying chunk size, chunk alignment, and memory allocation functions. A *chunk_size* or *alignment* of zero results in the default size or alignment respectively being used.

`int obstack_specify_allocation_with_arg (struct obstack *obstack-ptr, size_t chunk_size, size_t alignment, void *(*chunkfun) (void *, size_t), void (*freefun) (void *, void *), void *arg)` [Function]

Like `obstack_specify_allocation`, but specifying memory allocation functions that take an extra first argument, *arg*.

`obstack_alloc_failed_handler` [Variable]

The value of this variable is a pointer to a function that `obstack` uses when `obstack_chunk_alloc` fails to allocate memory. The default action is to print a message and abort. You should supply a function that either calls `exit` (see Section "Program Termination" in *The GNU C Library Reference Manual*) or `longjmp` and doesn't return.

```
void my_obstack_alloc_failed (void)
...
obstack_alloc_failed_handler = &my_obstack_alloc_failed;
```

2.3.1.3 Allocation in an Obstack

The most direct way to allocate an object in an obstack is with `obstack_alloc`, which is invoked almost like `malloc`.

`void * obstack_alloc (struct obstack *obstack-ptr, size_t size)` [Function]

This allocates an uninitialized block of *size* bytes in an obstack and returns its address. Here *obstack-ptr* specifies which obstack to allocate the block in; it is the address of the `struct obstack` object which represents the obstack. Each obstack macro requires you to specify an *obstack-ptr* as the first argument.

This macro calls the obstack's `obstack_chunk_alloc` function if it needs to allocate a new chunk of memory; it calls `obstack_alloc_failed_handler` if allocation of memory by `obstack_chunk_alloc` failed.

For example, here is a function that allocates a copy of a string *str* in a specific obstack, which is in the variable `string_obstack`:

```
struct obstack string_obstack;

char *
copystring (char *string)
{
  size_t len = strlen (string) + 1;
  char *s = (char *) obstack_alloc (&string_obstack, len);
  memcpy (s, string, len);
  return s;
}
```

To allocate a block with specified contents, use the macro `obstack_copy`.

`void * obstack_copy (struct obstack *obstack-ptr, void *address, size_t size)` [Function]

This allocates a block and initializes it by copying *size* bytes of data starting at *address*. It calls `obstack_alloc_failed_handler` if allocation of memory by `obstack_chunk_alloc` failed.

`void * obstack_copy0 (struct obstack *obstack-ptr, void *address, size_t size)` [Function]

Like `obstack_copy`, but appends an extra byte containing a null character. This extra byte is not counted in the argument *size*.

The `obstack_copy0` macro is convenient for copying a sequence of characters into an obstack as a null-terminated string. Here is an example of its use:

```
char *
obstack_savestring (char *addr, size_t size)
{
  return obstack_copy0 (&myobstack, addr, size);
}
```

Contrast this with the previous example of `savestring` using `malloc` (see Section "Basic Allocation" in *The GNU C Library Reference Manual*).

2.3.1.4 Freeing Objects in an Obstack

To free an object allocated in an obstack, use the macro `obstack_free`. Since the obstack is a stack of objects, freeing one object automatically frees all other objects allocated more recently in the same obstack.

`void obstack_free (struct obstack *obstack-ptr, void *object)` [Function]

If *object* is a null pointer, everything allocated in the obstack is freed. Otherwise, *object* must be the address of an object allocated in the obstack. Then *object* is freed, along with everything allocated in *obstack* since *object*.

Note that if *object* is a null pointer, the result is an uninitialized obstack. To free all memory in an obstack but leave it valid for further allocation, call `obstack_free` with the address of the first object allocated on the obstack:

```
obstack_free (obstack_ptr, first_object_allocated_ptr);
```

Recall that the objects in an obstack are grouped into chunks. When all the objects in a chunk become free, the obstack library automatically frees the chunk (see Section 2.3.1.2 [Preparing for Obstacks], page 3). Then other obstacks, or non-obstack allocation, can reuse the space of the chunk.

2.3.1.5 Obstack Functions and Macros

The interfaces for using obstacks are shown here as functions to specify the return type and argument types, but they are really defined as macros. This means that the arguments don't actually have types, but they generally behave as if they have the types shown. You can call these macros like functions, but you cannot use them in any other way (for example, you cannot take their address).

Calling the macros requires a special precaution: namely, the first operand (the obstack pointer) may not contain any side effects, because it may be computed more than once. For example, if you write this:

```
obstack_alloc (get_obstack (), 4);
```

you will find that `get_obstack` may be called several times. If you use `*obstack_list_ptr++` as the obstack pointer argument, you will get very strange results since the incrementation may occur several times.

If you use the GNU C compiler, this precaution is not necessary, because various language extensions in GNU C permit defining the macros so as to compute each argument only once.

Note that arguments other than the first will only be evaluated once, even when not using GNU C.

`obstack.h` does declare a number of functions, `_obstack_begin`, `_obstack_begin_1`, `_obstack_newchunk`, `_obstack_free`, and `_obstack_memory_used`. You should not call these directly.

2.3.1.6 Growing Objects

Because memory in obstack chunks is used sequentially, it is possible to build up an object step by step, adding one or more bytes at a time to the end of the object. With this technique, you do not need to know how much data you will put in the object until you come to the end of it. We call this the technique of *growing objects*. The special macros for adding data to the growing object are described in this section.

You don't need to do anything special when you start to grow an object. Using one of the macros to add data to the object automatically starts it. However, it is necessary to say explicitly when the object is finished. This is done with `obstack_finish`.

The actual address of the object thus built up is not known until the object is finished. Until then, it always remains possible that you will add so much data that the object must be copied into a new chunk.

While the obstack is in use for a growing object, you cannot use it for ordinary allocation of another object. If you try to do so, the space already added to the growing object will become part of the other object.

`void obstack_blank (struct obstack *`*`obstack-ptr`*`, size_t` *`size`*`)` [Function]

> The most basic macro for adding to a growing object is `obstack_blank`, which adds space without initializing it.

`void obstack_grow (struct obstack *`*`obstack-ptr`*`, void *`*`data`*`, size_t` *`size`*`)` [Function]

> To add a block of initialized space, use `obstack_grow`, which is the growing-object analogue of `obstack_copy`. It adds *size* bytes of data to the growing object, copying the contents from *data*.

`void obstack_grow0 (struct obstack *`*`obstack-ptr`*`, void *`*`data`*`, size_t` *`size`*`)` [Function]

> This is the growing-object analogue of `obstack_copy0`. It adds *size* bytes copied from *data*, followed by an additional null character.

`void obstack_1grow (struct obstack *`*`obstack-ptr`*`, char` *`c`*`)` [Function]

> To add one character at a time, use `obstack_1grow`. It adds a single byte containing *c* to the growing object.

`void obstack_ptr_grow (struct obstack *`*`obstack-ptr`*`, void *`*`data`*`)` [Function]

> Adding the value of a pointer one can use `obstack_ptr_grow`. It adds `sizeof (void *)` bytes containing the value of *data*.

`void obstack_int_grow (struct obstack *obstack-ptr, int data)` [Function]

A single value of type `int` can be added by using `obstack_int_grow`. It adds `sizeof (int)` bytes to the growing object and initializes them with the value of *data*.

`void * obstack_finish (struct obstack *obstack-ptr)` [Function]

When you are finished growing the object, use `obstack_finish` to close it off and return its final address.

Once you have finished the object, the obstack is available for ordinary allocation or for growing another object.

When you build an object by growing it, you will probably need to know afterward how long it became. You need not keep track of this as you grow the object, because you can find out the length from the obstack with `obstack_object_size`, before finishing the object.

`size_t obstack_object_size (struct obstack *obstack-ptr)` [Function]

This macro returns the current size of the growing object, in bytes. Remember to call `obstack_object_size` *before* finishing the object. After it is finished, `obstack_object_size` will return zero.

If you have started growing an object and wish to cancel it, you should finish it and then free it, like this:

```
obstack_free (obstack_ptr, obstack_finish (obstack_ptr));
```

This has no effect if no object was growing.

2.3.1.7 Extra Fast Growing Objects

The usual macros for growing objects incur overhead for checking whether there is room for the new growth in the current chunk. If you are frequently constructing objects in small steps of growth, this overhead can be significant.

You can reduce the overhead by using special "fast growth" macros that grow the object without checking. In order to have a robust program, you must do the checking yourself. If you do this checking in the simplest way each time you are about to add data to the object, you have not saved anything, because that is what the ordinary growth macros do. But if you can arrange to check less often, or check more efficiently, then you make the program faster.

`obstack_room` returns the amount of room available in the current chunk.

`size_t obstack_room (struct obstack *obstack-ptr)` [Function]

This returns the number of bytes that can be added safely to the current growing object (or to an object about to be started) in obstack *obstack* using the fast growth macros.

While you know there is room, you can use these fast growth macros for adding data to a growing object:

`void obstack_1grow_fast (struct obstack *obstack-ptr, char c)` [Function]

`obstack_1grow_fast` adds one byte containing the character *c* to the growing object in obstack *obstack-ptr*.

`void obstack_ptr_grow_fast (struct obstack *obstack-ptr, void *data)` [Function]

`obstack_ptr_grow_fast` adds `sizeof (void *)` bytes containing the value of *data* to the growing object in obstack *obstack-ptr*.

`void obstack_int_grow_fast (struct obstack *obstack-ptr, int data)` [Function]

`obstack_int_grow_fast` adds `sizeof (int)` bytes containing the value of *data* to the growing object in obstack *obstack-ptr*.

`void obstack_blank_fast (struct obstack *obstack-ptr, size_t size)` [Function]

`obstack_blank_fast` adds *size* bytes to the growing object in obstack *obstack-ptr* without initializing them.

When you check for space using `obstack_room` and there is not enough room for what you want to add, the fast growth macros are not safe. In this case, simply use the corresponding ordinary growth macro instead. Very soon this will copy the object to a new chunk; then there will be lots of room available again.

So, each time you use an ordinary growth macro, check afterward for sufficient space using `obstack_room`. Once the object is copied to a new chunk, there will be plenty of space again, so the program will start using the fast growth macros again.

Here is an example:

```
void
add_string (struct obstack *obstack, const char *ptr, size_t len)
{
  while (len > 0)
    {
      size_t room = obstack_room (obstack);
      if (room == 0)
        {
          /* Not enough room.  Add one character slowly,
             which may copy to a new chunk and make room.  */
          obstack_1grow (obstack, *ptr++);
          len--;
        }
      else
        {
          if (room > len)
            room = len;
          /* Add fast as much as we have room for. */
          len -= room;
          while (room-- > 0)
            obstack_1grow_fast (obstack, *ptr++);
        }
    }
}
```

You can use `obstack_blank_fast` with a "negative" size argument to make the current object smaller. Just don't try to shrink it beyond zero length—there's no telling what will happen if you do that. Earlier versions of obstacks allowed you to use `obstack_blank` to shrink objects. This will no longer work.

2.3.1.8 Status of an Obstack

Here are macros that provide information on the current status of allocation in an obstack. You can use them to learn about an object while still growing it.

`void * obstack_base (struct obstack *obstack-ptr)` [Function]

This macro returns the tentative address of the beginning of the currently growing object in *obstack-ptr*. If you finish the object immediately, it will have that address. If you make it larger first, it may outgrow the current chunk—then its address will change!

If no object is growing, this value says where the next object you allocate will start (once again assuming it fits in the current chunk).

`void * obstack_next_free (struct obstack *obstack-ptr)` [Function]

This macro returns the address of the first free byte in the current chunk of obstack *obstack-ptr*. This is the end of the currently growing object. If no object is growing, `obstack_next_free` returns the same value as `obstack_base`.

`size_t obstack_object_size (struct obstack *obstack-ptr)` [Function]

This macro returns the size in bytes of the currently growing object. This is equivalent to

```
((size_t) (obstack_next_free (obstack-ptr) - obstack_base (obstack-ptr)))
```

2.3.1.9 Alignment of Data in Obstacks

Each obstack has an *alignment boundary*; each object allocated in the obstack automatically starts on an address that is a multiple of the specified boundary. By default, this boundary is aligned so that the object can hold any type of data.

To access an obstack's alignment boundary, use the macro `obstack_alignment_mask`.

`size_t obstack_alignment_mask (struct obstack *obstack-ptr)` [Macro]

The value is a bit mask; a bit that is 1 indicates that the corresponding bit in the address of an object should be 0. The mask value should be one less than a power of 2; the effect is that all object addresses are multiples of that power of 2. The default value of the mask is a value that allows aligned objects to hold any type of data: for example, if its value is 3, any type of data can be stored at locations whose addresses are multiples of 4. A mask value of 0 means an object can start on any multiple of 1 (that is, no alignment is required).

The expansion of the macro `obstack_alignment_mask` is an lvalue, so you can alter the mask by assignment. For example, this statement:

```
obstack_alignment_mask (obstack_ptr) = 0;
```

has the effect of turning off alignment processing in the specified obstack.

Note that a change in alignment mask does not take effect until *after* the next time an object is allocated or finished in the obstack. If you are not growing an object, you can make the new alignment mask take effect immediately by calling `obstack_finish`. This will finish a zero-length object and then do proper alignment for the next object.

2.3.1.10 Obstack Chunks

Obstacks work by allocating space for themselves in large chunks, and then parceling out space in the chunks to satisfy your requests. Chunks are normally 4096 bytes long unless you specify a different chunk size. The chunk size includes 8 bytes of overhead that are not actually used for storing objects. Regardless of the specified size, longer chunks will be allocated when necessary for long objects.

The obstack library allocates chunks by calling the function `obstack_chunk_alloc`, which you must define. When a chunk is no longer needed because you have freed all the objects in it, the obstack library frees the chunk by calling `obstack_chunk_free`, which you must also define.

These two must be defined (as macros) or declared (as functions) in each source file that uses `obstack_init` (see Section 2.3.1.1 [Creating Obstacks], page 3). Most often they are defined as macros like this:

```
#define obstack_chunk_alloc malloc
#define obstack_chunk_free free
```

Note that these are simple macros (no arguments). Macro definitions with arguments will not work! It is necessary that `obstack_chunk_alloc` or `obstack_chunk_free`, alone, expand into a function name if it is not itself a function name.

If you allocate chunks with `malloc`, the chunk size should be a power of 2. The default chunk size, 4096, was chosen because it is long enough to satisfy many typical requests on the obstack yet short enough not to waste too much memory in the portion of the last chunk not yet used.

`size_t obstack_chunk_size (struct obstack *`*`obstack-ptr`*`)` [Macro]

This returns the chunk size of the given obstack.

Since this macro expands to an lvalue, you can specify a new chunk size by assigning it a new value. Doing so does not affect the chunks already allocated, but will change the size of chunks allocated for that particular obstack in the future. It is unlikely to be useful to make the chunk size smaller, but making it larger might improve efficiency if you are allocating many objects whose size is comparable to the chunk size. Here is how to do so cleanly:

```
if (obstack_chunk_size (obstack_ptr) < new-chunk-size)
  obstack_chunk_size (obstack_ptr) = new-chunk-size;
```

2.3.1.11 Summary of Obstack Macros

Here is a summary of all the macros associated with obstacks. Each takes the address of an obstack (`struct obstack *`) as its first argument.

`int obstack_init (struct obstack *`*`obstack-ptr`*`)`

Initialize use of an obstack. See Section 2.3.1.1 [Creating Obstacks], page 3.

`int obstack_begin (struct obstack *`*`obstack-ptr`*`, size_t chunk_size)`

Initialize use of an obstack, with an initial chunk of *chunk_size* bytes.

`int obstack_specify_allocation (struct obstack *obstack-ptr, size_t chunk_size, size_t alignment, void *(*chunkfun) (size_t), void (*freefun) (void *))`

Initialize use of an obstack, specifying intial chunk size, chunk alignment, and memory allocation functions.

`int obstack_specify_allocation_with_arg (struct obstack *obstack-ptr, size_t chunk_size, size_t alignment, void *(*chunkfun) (void *, size_t), void (*freefun) (void *, void *), void *arg)`

Like `obstack_specify_allocation`, but specifying memory allocation functions that take an extra first argument, *arg*.

`void *obstack_alloc (struct obstack *obstack-ptr, size_t size)`

Allocate an object of *size* uninitialized bytes. See Section 2.3.1.3 [Allocation in an Obstack], page 5.

`void *obstack_copy (struct obstack *obstack-ptr, void *address, size_t size)`

Allocate an object of *size* bytes, with contents copied from *address*. See Section 2.3.1.3 [Allocation in an Obstack], page 5.

`void *obstack_copy0 (struct obstack *obstack-ptr, void *address, size_t size)`

Allocate an object of *size*+1 bytes, with *size* of them copied from *address*, followed by a null character at the end. See Section 2.3.1.3 [Allocation in an Obstack], page 5.

`void obstack_free (struct obstack *obstack-ptr, void *object)`

Free *object* (and everything allocated in the specified obstack more recently than *object*). See Section 2.3.1.4 [Freeing Obstack Objects], page 6.

`void obstack_blank (struct obstack *obstack-ptr, size_t size)`

Add *size* uninitialized bytes to a growing object. See Section 2.3.1.6 [Growing Objects], page 7.

`void obstack_grow (struct obstack *obstack-ptr, void *address, size_t size)`

Add *size* bytes, copied from *address*, to a growing object. See Section 2.3.1.6 [Growing Objects], page 7.

`void obstack_grow0 (struct obstack *obstack-ptr, void *address, size_t size)`

Add *size* bytes, copied from *address*, to a growing object, and then add another byte containing a null character. See Section 2.3.1.6 [Growing Objects], page 7.

`void obstack_1grow (struct obstack *obstack-ptr, char data-char)`

Add one byte containing *data-char* to a growing object. See Section 2.3.1.6 [Growing Objects], page 7.

`void *obstack_finish (struct obstack *obstack-ptr)`

Finalize the object that is growing and return its permanent address. See Section 2.3.1.6 [Growing Objects], page 7.

`size_t obstack_object_size (struct obstack *obstack-ptr)`

Get the current size of the currently growing object. See Section 2.3.1.6 [Growing Objects], page 7.

`void obstack_blank_fast (struct obstack *`*`obstack-ptr`*`, size_t `*`size`*`)`
: Add *size* uninitialized bytes to a growing object without checking that there is enough room. See Section 2.3.1.7 [Extra Fast Growing], page 8.

`void obstack_1grow_fast (struct obstack *`*`obstack-ptr`*`, char `*`data-char`*`)`
: Add one byte containing *data-char* to a growing object without checking that there is enough room. See Section 2.3.1.7 [Extra Fast Growing], page 8.

`size_t obstack_room (struct obstack *`*`obstack-ptr`*`)`
: Get the amount of room now available for growing the current object. See Section 2.3.1.7 [Extra Fast Growing], page 8.

`size_t obstack_alignment_mask (struct obstack *`*`obstack-ptr`*`)`
: The mask used for aligning the beginning of an object. This is an lvalue. See Section 2.3.1.9 [Obstacks Data Alignment], page 10.

`size_t obstack_chunk_size (struct obstack *`*`obstack-ptr`*`)`
: The size for allocating chunks. This is an lvalue. See Section 2.3.1.10 [Obstack Chunks], page 11.

`void *obstack_base (struct obstack *`*`obstack-ptr`*`)`
: Tentative starting address of the currently growing object. See Section 2.3.1.8 [Status of an Obstack], page 10.

`void *obstack_next_free (struct obstack *`*`obstack-ptr`*`)`
: Address just after the end of the currently growing object. See Section 2.3.1.8 [Status of an Obstack], page 10.

3 Function, Variable, and Macro Listing.

`void* alloca (size_t size)` [Replacement]

This function allocates memory which will be automatically reclaimed after the procedure exits. The `libiberty` implementation does not free the memory immediately but will do so eventually during subsequent calls to this function. Memory is allocated using `xmalloc` under normal circumstances.

The header file `alloca-conf.h` can be used in conjunction with the GNU Autoconf test `AC_FUNC_ALLOCA` to test for and properly make available this function. The `AC_FUNC_ALLOCA` test requires that client code use a block of preprocessor code to be safe (see the Autoconf manual for more); this header incorporates that logic and more, including the possibility of a GCC built-in function.

`int asprintf (char **resptr, const char *format, ...)` [Extension]

Like `sprintf`, but instead of passing a pointer to a buffer, you pass a pointer to a pointer. This function will compute the size of the buffer needed, allocate memory with `malloc`, and store a pointer to the allocated memory in `*resptr`. The value returned is the same as `sprintf` would return. If memory could not be allocated, minus one is returned and `NULL` is stored in `*resptr`.

`int atexit (void (*f)())` [Supplemental]

Causes function *f* to be called at exit. Returns 0.

`char* basename (const char *name)` [Supplemental]

Returns a pointer to the last component of pathname *name*. Behavior is undefined if the pathname ends in a directory separator.

`int bcmp (char *x, char *y, int count)` [Supplemental]

Compares the first *count* bytes of two areas of memory. Returns zero if they are the same, nonzero otherwise. Returns zero if *count* is zero. A nonzero result only indicates a difference, it does not indicate any sorting order (say, by having a positive result mean *x* sorts before *y*).

`void bcopy (char *in, char *out, int length)` [Supplemental]

Copies *length* bytes from memory region *in* to region *out*. The use of `bcopy` is deprecated in new programs.

`void* bsearch (const void *key, const void *base, size_t nmemb, size_t size, int (*compar)(const void *, const void *))` [Supplemental]

Performs a search over an array of *nmemb* elements pointed to by *base* for a member that matches the object pointed to by *key*. The size of each member is specified by *size*. The array contents should be sorted in ascending order according to the *compar* comparison function. This routine should take two arguments pointing to the *key* and to an array member, in that order, and should return an integer less than, equal to, or greater than zero if the *key* object is respectively less than, matching, or greater than the array member.

`void* bsearch_r (const void *key, const void *base, size_t nmemb, size_t size, int (*compar)(const void *, const void *, void *), void *arg)` [Supplemental]

Performs a search over an array of *nmemb* elements pointed to by *base* for a member that matches the object pointed to by *key*. The size of each member is specified by *size*. The array contents should be sorted in ascending order according to the *compar* comparison function. This routine should take three arguments: the first two point to the *key* and to an array member, and the last is passed down unchanged from `bsearch_r`'s last argument. It should return an integer less than, equal to, or greater than zero if the *key* object is respectively less than, matching, or greater than the array member.

`char** buildargv (char *sp)` [Extension]

Given a pointer to a string, parse the string extracting fields separated by whitespace and optionally enclosed within either single or double quotes (which are stripped off), and build a vector of pointers to copies of the string for each field. The input string remains unchanged. The last element of the vector is followed by a `NULL` element.

All of the memory for the pointer array and copies of the string is obtained from `xmalloc`. All of the memory can be returned to the system with the single function call `freeargv`, which takes the returned result of `buildargv`, as it's argument.

Returns a pointer to the argument vector if successful. Returns `NULL` if *sp* is `NULL` or if there is insufficient memory to complete building the argument vector.

If the input is a null string (as opposed to a `NULL` pointer), then buildarg returns an argument vector that has one arg, a null string.

`void bzero (char *mem, int count)` [Supplemental]

Zeros *count* bytes starting at *mem*. Use of this function is deprecated in favor of `memset`.

`void* calloc (size_t nelem, size_t elsize)` [Supplemental]

Uses `malloc` to allocate storage for *nelem* objects of *elsize* bytes each, then zeros the memory.

`int canonical_filename_eq (const char *a, const char *b)` [Extension]

Return non-zero if file names *a* and *b* are equivalent. This function compares the canonical versions of the filenames as returned by `lrealpath()`, so that so that different file names pointing to the same underlying file are treated as being identical.

`char* choose_temp_base (void)` [Extension]

Return a prefix for temporary file names or `NULL` if unable to find one. The current directory is chosen if all else fails so the program is exited if a temporary directory can't be found (`mktemp` fails). The buffer for the result is obtained with `xmalloc`.

This function is provided for backwards compatibility only. Its use is not recommended.

`const char* choose_tmpdir ()` [Replacement]

Returns a pointer to a directory path suitable for creating temporary files in.

`long clock (void)` [Supplemental]

Returns an approximation of the CPU time used by the process as a `clock_t`; divide this number by '`CLOCKS_PER_SEC`' to get the number of seconds used.

`char* concat (const char *s1, const char *s2, ..., NULL)` [Extension]

Concatenate zero or more of strings and return the result in freshly `xmalloc`ed memory. The argument list is terminated by the first `NULL` pointer encountered. Pointers to empty strings are ignored.

`int countargv (char * const *argv)` [Extension]

Return the number of elements in *argv*. Returns zero if *argv* is NULL.

`unsigned int crc32 (const unsigned char *buf, int len, unsigned int init)` [Extension]

Compute the 32-bit CRC of *buf* which has length *len*. The starting value is *init*; this may be used to compute the CRC of data split across multiple buffers by passing the return value of each call as the *init* parameter of the next.

This is used by the `gdb` remote protocol for the '`qCRC`' command. In order to get the same results as gdb for a block of data, you must pass the first CRC parameter as `0xffffffff`.

This CRC can be specified as:

Width : 32 Poly : 0x04c11db7 Init : parameter, typically 0xffffffff RefIn : false RefOut : false XorOut : 0

This differs from the "standard" CRC-32 algorithm in that the values are not reflected, and there is no final XOR value. These differences make it easy to compose the values of multiple blocks.

`char** dupargv (char * const *vector)` [Extension]

Duplicate an argument vector. Simply scans through *vector*, duplicating each argument until the terminating `NULL` is found. Returns a pointer to the argument vector if successful. Returns `NULL` if there is insufficient memory to complete building the argument vector.

`int errno_max (void)` [Extension]

Returns the maximum `errno` value for which a corresponding symbolic name or message is available. Note that in the case where we use the `sys_errlist` supplied by the system, it is possible for there to be more symbolic names than messages, or vice versa. In fact, the manual page for `perror(3C)` explicitly warns that one should check the size of the table (`sys_nerr`) before indexing it, since new error codes may be added to the system before they are added to the table. Thus `sys_nerr` might be smaller than value implied by the largest `errno` value defined in `<errno.h>`.

We return the maximum value that can be used to obtain a meaningful symbolic name or message.

`void expandargv (int *argcp, char ***argvp)` [Extension]

The *argcp* and `argvp` arguments are pointers to the usual `argc` and `argv` arguments to `main`. This function looks for arguments that begin with the character '`@`'. Any such arguments are interpreted as "response files". The contents of the response file are interpreted as additional command line options. In particular, the file is separated into whitespace-separated strings; each such string is taken as a command-line option. The new options are inserted in place of the option naming the response file, and `*argcp` and `*argvp` will be updated. If the value of `*argvp` is modified by this function, then the new value has been dynamically allocated and can be deallocated by the caller with `freeargv`. However, most callers will simply call `expandargv` near the beginning of `main` and allow the operating system to free the memory when the program exits.

`int fdmatch (int fd1, int fd2)` [Extension]

Check to see if two open file descriptors refer to the same file. This is useful, for example, when we have an open file descriptor for an unnamed file, and the name of a file that we believe to correspond to that fd. This can happen when we are exec'd with an already open file (`stdout` for example) or from the SVR4 `/proc` calls that return open file descriptors for mapped address spaces. All we have to do is open the file by name and check the two file descriptors for a match, which is done by comparing major and minor device numbers and inode numbers.

`FILE * fdopen_unlocked (int fildes, const char * mode)` [Extension]

Opens and returns a `FILE` pointer via `fdopen`. If the operating system supports it, ensure that the stream is setup to avoid any multi-threaded locking. Otherwise return the `FILE` pointer unchanged.

`int ffs (int valu)` [Supplemental]

Find the first (least significant) bit set in *valu*. Bits are numbered from right to left, starting with bit 1 (corresponding to the value 1). If *valu* is zero, zero is returned.

`int filename_cmp (const char *s1, const char *s2)` [Extension]

Return zero if the two file names *s1* and *s2* are equivalent. If not equivalent, the returned value is similar to what `strcmp` would return. In other words, it returns a negative value if *s1* is less than *s2*, or a positive value if *s2* is greater than *s2*.

This function does not normalize file names. As a result, this function will treat filenames that are spelled differently as different even in the case when the two filenames point to the same underlying file. However, it does handle the fact that on DOS-like file systems, forward and backward slashes are equal.

`int filename_eq (const void *s1, const void *s2)` [Extension]

Return non-zero if file names *s1* and *s2* are equivalent. This function is for use with hashtab.c hash tables.

`hashval_t filename_hash (const void *s)` [Extension]

Return the hash value for file name *s* that will be compared using filename_cmp. This function is for use with hashtab.c hash tables.

`int filename_ncmp (const char *s1, const char *s2, size_t n)` [Extension]

Return zero if the two file names *s1* and *s2* are equivalent in range *n*. If not equivalent, the returned value is similar to what `strncmp` would return. In other words, it returns a negative value if *s1* is less than *s2*, or a positive value if *s2* is greater than *s2*.

This function does not normalize file names. As a result, this function will treat filenames that are spelled differently as different even in the case when the two filenames point to the same underlying file. However, it does handle the fact that on DOS-like file systems, forward and backward slashes are equal.

`int fnmatch (const char *pattern, const char *string, int flags)` [Replacement]

Matches *string* against *pattern*, returning zero if it matches, `FNM_NOMATCH` if not. *pattern* may contain the wildcards `?` to match any one character, `*` to match any zero or more characters, or a set of alternate characters in square brackets, like '`[a-gt8]`', which match one character (`a` through `g`, or `t`, or `8`, in this example) if that one character is in the set. A set may be inverted (i.e., match anything except what's in the set) by giving `^` or `!` as the first character in the set. To include those characters in the set, list them as anything other than the first character of the set. To include a dash in the set, list it last in the set. A backslash character makes the following character not special, so for example you could match against a literal asterisk with '`\*`'. To match a literal backslash, use '`\\`'.

`flags` controls various aspects of the matching process, and is a boolean OR of zero or more of the following values (defined in `<fnmatch.h>`):

`FNM_PATHNAME`
`FNM_FILE_NAME`
: *string* is assumed to be a path name. No wildcard will ever match `/`.

`FNM_NOESCAPE`
: Do not interpret backslashes as quoting the following special character.

`FNM_PERIOD`
: A leading period (at the beginning of *string*, or if `FNM_PATHNAME` after a slash) is not matched by `*` or `?` but must be matched explicitly.

`FNM_LEADING_DIR`
: Means that *string* also matches *pattern* if some initial part of *string* matches, and is followed by `/` and zero or more characters. For example, '`foo*`' would match either '`foobar`' or '`foobar/grill`'.

`FNM_CASEFOLD`
: Ignores case when performing the comparison.

`FILE * fopen_unlocked (const char *path, const char * mode)` [Extension]

Opens and returns a `FILE` pointer via `fopen`. If the operating system supports it, ensure that the stream is setup to avoid any multi-threaded locking. Otherwise return the `FILE` pointer unchanged.

`void freeargv (char **vector)` [Extension]

Free an argument vector that was built using `buildargv`. Simply scans through *vector*, freeing the memory for each argument until the terminating `NULL` is found, and then frees *vector* itself.

`FILE * freopen_unlocked (const char * path, const char * mode, FILE * stream)` [Extension]

Opens and returns a `FILE` pointer via `freopen`. If the operating system supports it, ensure that the stream is setup to avoid any multi-threaded locking. Otherwise return the `FILE` pointer unchanged.

`long get_run_time (void)` [Replacement]

Returns the time used so far, in microseconds. If possible, this is the time used by this process, else it is the elapsed time since the process started.

`char* getcwd (char *pathname, int len)` [Supplemental]

Copy the absolute pathname for the current working directory into *pathname*, which is assumed to point to a buffer of at least *len* bytes, and return a pointer to the buffer. If the current directory's path doesn't fit in *len* characters, the result is `NULL` and `errno` is set. If *pathname* is a null pointer, `getcwd` will obtain *len* bytes of space using `malloc`.

`int getpagesize (void)` [Supplemental]

Returns the number of bytes in a page of memory. This is the granularity of many of the system memory management routines. No guarantee is made as to whether or not it is the same as the basic memory management hardware page size.

`char* getpwd (void)` [Supplemental]

Returns the current working directory. This implementation caches the result on the assumption that the process will not call `chdir` between calls to `getpwd`.

`int gettimeofday (struct timeval *tp, void *tz)` [Supplemental]

Writes the current time to *tp*. This implementation requires that *tz* be NULL. Returns 0 on success, -1 on failure.

`void hex_init (void)` [Extension]

Initializes the array mapping the current character set to corresponding hex values. This function must be called before any call to `hex_p` or `hex_value`. If you fail to call it, a default ASCII-based table will normally be used on ASCII systems.

`int hex_p (int c)` [Extension]

Evaluates to non-zero if the given character is a valid hex character, or zero if it is not. Note that the value you pass will be cast to `unsigned char` within the macro.

`unsigned int hex_value (int c)` [Extension]

Returns the numeric equivalent of the given character when interpreted as a hexadecimal digit. The result is undefined if you pass an invalid hex digit. Note that the value you pass will be cast to `unsigned char` within the macro.

The `hex_value` macro returns `unsigned int`, rather than signed `int`, to make it easier to use in parsing addresses from hex dump files: a signed `int` would be sign-extended when converted to a wider unsigned type — like `bfd_vma`, on some systems.

`HOST_CHARSET` [Extension]

This macro indicates the basic character set and encoding used by the host: more precisely, the encoding used for character constants in preprocessor '`#if`' statements (the C "execution character set"). It is defined by `safe-ctype.h`, and will be an integer constant with one of the following values:

`HOST_CHARSET_UNKNOWN`
: The host character set is unknown - that is, not one of the next two possibilities.

`HOST_CHARSET_ASCII`
: The host character set is ASCII.

`HOST_CHARSET_EBCDIC`
: The host character set is some variant of EBCDIC. (Only one of the nineteen EBCDIC varying characters is tested; exercise caution.)

`htab_t htab_create_typed_alloc (size_t` *size*`, htab_hash` *hash_f*`, htab_eq` *eq_f*`, htab_del` *del_f*`, htab_alloc` *alloc_tab_f*`, htab_alloc` *alloc_f*`, htab_free` *free_f*`)` [Supplemental]

This function creates a hash table that uses two different allocators *alloc_tab_f* and *alloc_f* to use for allocating the table itself and its entries respectively. This is useful when variables of different types need to be allocated with different allocators.

The created hash table is slightly larger than *size* and it is initially empty (all the hash table entries are `HTAB_EMPTY_ENTRY`). The function returns the created hash table, or `NULL` if memory allocation fails.

`char* index (char *`*s*`, int` *c*`)` [Supplemental]

Returns a pointer to the first occurrence of the character *c* in the string *s*, or `NULL` if not found. The use of `index` is deprecated in new programs in favor of `strchr`.

`void insque (struct qelem *`*elem*`, struct qelem *`*pred*`)` [Supplemental]
`void remque (struct qelem *`*elem*`)` [Supplemental]

Routines to manipulate queues built from doubly linked lists. The `insque` routine inserts *elem* in the queue immediately after *pred*. The `remque` routine removes *elem* from its containing queue. These routines expect to be passed pointers to structures which have as their first members a forward pointer and a back pointer, like this prototype (although no prototype is provided):

```
struct qelem {
  struct qelem *q_forw;
  struct qelem *q_back;
  char q_data[];
};
```

`ISALPHA (`*c*`)` [Extension]
`ISALNUM (`*c*`)` [Extension]
`ISBLANK (`*c*`)` [Extension]
`ISCNTRL (`*c*`)` [Extension]
`ISDIGIT (`*c*`)` [Extension]
`ISGRAPH (`*c*`)` [Extension]
`ISLOWER (`*c*`)` [Extension]
`ISPRINT (`*c*`)` [Extension]

`ISPUNCT` (*c*) [Extension]
`ISSPACE` (*c*) [Extension]
`ISUPPER` (*c*) [Extension]
`ISXDIGIT` (*c*) [Extension]

These twelve macros are defined by `safe-ctype.h`. Each has the same meaning as the corresponding macro (with name in lowercase) defined by the standard header `ctype.h`. For example, `ISALPHA` returns true for alphabetic characters and false for others. However, there are two differences between these macros and those provided by `ctype.h`:

- These macros are guaranteed to have well-defined behavior for all values representable by `signed char` and `unsigned char`, and for `EOF`.
- These macros ignore the current locale; they are true for these fixed sets of characters:

ALPHA	*A-Za-z*
ALNUM	*A-Za-z0-9*
BLANK	*space tab*
CNTRL	!PRINT
DIGIT	*0-9*
GRAPH	ALNUM || PUNCT
LOWER	*a-z*
PRINT	GRAPH || *space*
PUNCT	*`~!@#$%^&*()_-=+[{]}\\|;:'",<.>/?*
SPACE	*space tab \n \r \f \v*
UPPER	*A-Z*
XDIGIT	*0-9A-Fa-f*

Note that, if the host character set is ASCII or a superset thereof, all these macros will return false for all values of `char` outside the range of 7-bit ASCII. In particular, both ISPRINT and ISCNTRL return false for characters with numeric values from 128 to 255.

`ISIDNUM` (*c*) [Extension]
`ISIDST` (*c*) [Extension]
`IS_VSPACE` (*c*) [Extension]
`IS_NVSPACE` (*c*) [Extension]
`IS_SPACE_OR_NUL` (*c*) [Extension]
`IS_ISOBASIC` (*c*) [Extension]

These six macros are defined by `safe-ctype.h` and provide additional character classes which are useful when doing lexical analysis of C or similar languages. They are true for the following sets of characters:

IDNUM	*A-Za-z0-9_*
IDST	*A-Za-z_*
VSPACE	*\r \n*
NVSPACE	*space tab \f \v \0*
SPACE_OR_NUL	VSPACE || NVSPACE
ISOBASIC	VSPACE || NVSPACE || PRINT

`const char* lbasename (const char *name)` [Replacement]

Given a pointer to a string containing a typical pathname ('`/usr/src/cmd/ls/ls.c`' for example), returns a pointer to the last component of the pathname ('`ls.c`' in this case). The returned pointer is guaranteed to lie within the original string. This latter fact is not true of many vendor C libraries, which return special strings or modify the passed strings for particular input.

In particular, the empty string returns the same empty string, and a path ending in `/` returns the empty string after it.

`const char* lrealpath (const char *name)` [Replacement]

Given a pointer to a string containing a pathname, returns a canonical version of the filename. Symlinks will be resolved, and "." and ".." components will be simplified. The returned value will be allocated using `malloc`, or `NULL` will be returned on a memory allocation error.

`const char* make_relative_prefix (const char *progname, const char *bin_prefix, const char *prefix)` [Extension]

Given three paths *progname*, *bin_prefix*, *prefix*, return the path that is in the same position relative to *progname*'s directory as *prefix* is relative to *bin_prefix*. That is, a string starting with the directory portion of *progname*, followed by a relative pathname of the difference between *bin_prefix* and *prefix*.

If *progname* does not contain any directory separators, `make_relative_prefix` will search `PATH` to find a program named *progname*. Also, if *progname* is a symbolic link, the symbolic link will be resolved.

For example, if *bin_prefix* is `/alpha/beta/gamma/gcc/delta`, *prefix* is `/alpha/beta/gamma/omega/`, and *progname* is `/red/green/blue/gcc`, then this function will return `/red/green/blue/../../omega/`.

The return value is normally allocated via `malloc`. If no relative prefix can be found, return `NULL`.

`char* make_temp_file (const char *suffix)` [Replacement]

Return a temporary file name (as a string) or `NULL` if unable to create one. *suffix* is a suffix to append to the file name. The string is `malloc`ed, and the temporary file has been created.

`void* memchr (const void *s, int c, size_t n)` [Supplemental]

This function searches memory starting at `*s` for the character *c*. The search only ends with the first occurrence of *c*, or after *length* characters; in particular, a null character does not terminate the search. If the character *c* is found within *length* characters of `*s`, a pointer to the character is returned. If *c* is not found, then `NULL` is returned.

`void* memrchr (const void *s, int c, size_t n)` [Supplemental]

This function searches memory for the character *c* in reverse order, starting at `*s+n-1` . The search only ends with the first occurrence of *c*, or when the start us reached; in particular, a null character does not terminate the search. If the character *c* is found within *length* characters of `*s`, a pointer to the character is returned. If *c* is not found, then `NULL` is returned.

`int memcmp (const void *x, const void *y, size_t count)` [Supplemental]

Compares the first *count* bytes of two areas of memory. Returns zero if they are the same, a value less than zero if *x* is lexically less than *y*, or a value greater than zero if *x* is lexically greater than *y*. Note that lexical order is determined as if comparing unsigned char arrays.

`void* memcpy (void *out, const void *in, size_t length)` [Supplemental]

Copies *length* bytes from memory region *in* to region *out*. Returns a pointer to *out*.

`void* memmem (const void *haystack, size_t haystack_len const void *needle, size_t needle_len)` [Supplemental]

Returns a pointer to the first occurrence of *needle* (length *needle_len*) in *haystack* (length *haystack_len*). Returns `NULL` if not found.

`void* memmove (void *from, const void *to, size_t count)` [Supplemental]

Copies *count* bytes from memory area *from* to memory area *to*, returning a pointer to *to*.

`void* mempcpy (void *out, const void *in, size_t length)` [Supplemental]

Copies *length* bytes from memory region *in* to region *out*. Returns a pointer to *out* + *length*.

`void* memset (void *s, int c, size_t count)` [Supplemental]

Sets the first *count* bytes of *s* to the constant byte *c*, returning a pointer to *s*.

`int mkstemps (char *pattern, int suffix_len)` [Replacement]

Generate a unique temporary file name from *pattern*. *pattern* has the form:

```
path/ccXXXXXXsuffix
```

suffix_len tells us how long *suffix* is (it can be zero length). The last six characters of *pattern* before *suffix* must be '`XXXXXX`'; they are replaced with a string that makes the filename unique. Returns a file descriptor open on the file for reading and writing.

`void pex_free (struct pex_obj obj)` [Extension]

Clean up and free all data associated with *obj*. If you have not yet called `pex_get_times` or `pex_get_status`, this will try to kill the subprocesses.

`int pex_get_status (struct pex_obj *obj, int count, int *vector)` [Extension]

Returns the exit status of all programs run using *obj*. *count* is the number of results expected. The results will be placed into *vector*. The results are in the order of the calls to `pex_run`. Returns 0 on error, 1 on success.

`int pex_get_times (struct pex_obj *obj, int count, struct pex_time *vector)` [Extension]

Returns the process execution times of all programs run using *obj*. *count* is the number of results expected. The results will be placed into *vector*. The results are in the order of the calls to `pex_run`. Returns 0 on error, 1 on success.

`struct pex_time` has the following fields of the type `unsigned long`: `user_seconds`, `user_microseconds`, `system_seconds`, `system_microseconds`. On systems which do not support reporting process times, all the fields will be set to `0`.

`struct pex_obj * pex_init (int` *flags*`, const char *`*pname*`, const char *`*tempbase*`)` [Extension]

Prepare to execute one or more programs, with standard output of each program fed to standard input of the next. This is a system independent interface to execute a pipeline.

flags is a bitwise combination of the following:

`PEX_RECORD_TIMES`
: Record subprocess times if possible.

`PEX_USE_PIPES`
: Use pipes for communication between processes, if possible.

`PEX_SAVE_TEMPS`
: Don't delete temporary files used for communication between processes.

pname is the name of program to be executed, used in error messages. *tempbase* is a base name to use for any required temporary files; it may be `NULL` to use a randomly chosen name.

`FILE * pex_input_file (struct pex_obj *`*obj*`, int` *flags*`, const char *`*in_name*`)` [Extension]

Return a stream for a temporary file to pass to the first program in the pipeline as input.

The name of the input file is chosen according to the same rules `pex_run` uses to choose output file names, based on *in_name*, *obj* and the `PEX_SUFFIX` bit in *flags*.

Don't call `fclose` on the returned stream; the first call to `pex_run` closes it automatically.

If *flags* includes `PEX_BINARY_OUTPUT`, open the stream in binary mode; otherwise, open it in the default mode. Including `PEX_BINARY_OUTPUT` in *flags* has no effect on Unix.

`FILE * pex_input_pipe (struct pex_obj *`*obj*`, int` *binary*`)` [Extension]

Return a stream *fp* for a pipe connected to the standard input of the first program in the pipeline; *fp* is opened for writing. You must have passed `PEX_USE_PIPES` to the `pex_init` call that returned *obj*.

You must close *fp* using `fclose` yourself when you have finished writing data to the pipeline.

The file descriptor underlying *fp* is marked not to be inherited by child processes.

On systems that do not support pipes, this function returns `NULL`, and sets `errno` to `EINVAL`. If you would like to write code that is portable to all systems the `pex` functions support, consider using `pex_input_file` instead.

There are two opportunities for deadlock using `pex_input_pipe`:

- Most systems' pipes can buffer only a fixed amount of data; a process that writes to a full pipe blocks. Thus, if you write to `fp` before starting the first process, you run the risk of blocking when there is no child process yet to read the data and allow you to continue. `pex_input_pipe` makes no promises about the size of the pipe's buffer, so if you need to write any data at all before starting the first process in the pipeline, consider using `pex_input_file` instead.

- Using `pex_input_pipe` and `pex_read_output` together may also cause deadlock. If the output pipe fills up, so that each program in the pipeline is waiting for the next to read more data, and you fill the input pipe by writing more data to *fp*, then there is no way to make progress: the only process that could read data from the output pipe is you, but you are blocked on the input pipe.

`const char * pex_one (int *flags*, const char **executable*, char * const **argv*, const char **pname*, const char **outname*, const char **errname*, int **status*, int **err*)` [Extension]

An interface to permit the easy execution of a single program. The return value and most of the parameters are as for a call to `pex_run`. *flags* is restricted to a combination of `PEX_SEARCH`, `PEX_STDERR_TO_STDOUT`, and `PEX_BINARY_OUTPUT`. *outname* is interpreted as if `PEX_LAST` were set. On a successful return, `*status` will be set to the exit status of the program.

`FILE * pex_read_err (struct pex_obj **obj*, int *binary*)` [Extension]

Returns a `FILE` pointer which may be used to read the standard error of the last program in the pipeline. When this is used, `PEX_LAST` should not be used in a call to `pex_run`. After this is called, `pex_run` may no longer be called with the same *obj*. *binary* should be non-zero if the file should be opened in binary mode. Don't call `fclose` on the returned file; it will be closed by `pex_free`.

`FILE * pex_read_output (struct pex_obj **obj*, int *binary*)` [Extension]

Returns a `FILE` pointer which may be used to read the standard output of the last program in the pipeline. When this is used, `PEX_LAST` should not be used in a call to `pex_run`. After this is called, `pex_run` may no longer be called with the same *obj*. *binary* should be non-zero if the file should be opened in binary mode. Don't call `fclose` on the returned file; it will be closed by `pex_free`.

`const char * pex_run (struct pex_obj **obj*, int *flags*, const char **executable*, char * const **argv*, const char **outname*, const char **errname*, int **err*)` [Extension]

Execute one program in a pipeline. On success this returns `NULL`. On failure it returns an error message, a statically allocated string.

obj is returned by a previous call to `pex_init`.

flags is a bitwise combination of the following:

`PEX_LAST` This must be set on the last program in the pipeline. In particular, it should be set when executing a single program. The standard output of the program will be sent to *outname*, or, if *outname* is `NULL`, to the standard output of the calling program. Do *not* set this bit if you want to call `pex_read_output` (described below). After a call to `pex_run` with this bit set, *pex_run* may no longer be called with the same *obj*.

`PEX_SEARCH`
: Search for the program using the user's executable search path.

`PEX_SUFFIX`
: *outname* is a suffix. See the description of *outname*, below.

`PEX_STDERR_TO_STDOUT`
: Send the program's standard error to standard output, if possible.

`PEX_BINARY_INPUT`
`PEX_BINARY_OUTPUT`
`PEX_BINARY_ERROR`
: The standard input (output or error) of the program should be read (written) in binary mode rather than text mode. These flags are ignored on systems which do not distinguish binary mode and text mode, such as Unix. For proper behavior these flags should match appropriately—a call to `pex_run` using `PEX_BINARY_OUTPUT` should be followed by a call using `PEX_BINARY_INPUT`.

`PEX_STDERR_TO_PIPE`
: Send the program's standard error to a pipe, if possible. This flag cannot be specified together with `PEX_STDERR_TO_STDOUT`. This flag can be specified only on the last program in pipeline.

executable is the program to execute. *argv* is the set of arguments to pass to the program; normally `argv[0]` will be a copy of *executable*.

outname is used to set the name of the file to use for standard output. There are two cases in which no output file will be used:

1. if `PEX_LAST` is not set in *flags*, and `PEX_USE_PIPES` was set in the call to `pex_init`, and the system supports pipes
2. if `PEX_LAST` is set in *flags*, and *outname* is `NULL`

Otherwise the code will use a file to hold standard output. If `PEX_LAST` is not set, this file is considered to be a temporary file, and it will be removed when no longer needed, unless `PEX_SAVE_TEMPS` was set in the call to `pex_init`.

There are two cases to consider when setting the name of the file to hold standard output.

1. `PEX_SUFFIX` is set in *flags*. In this case *outname* may not be `NULL`. If the *tempbase* parameter to `pex_init` was not `NULL`, then the output file name is the concatenation of *tempbase* and *outname*. If *tempbase* was `NULL`, then the output file name is a random file name ending in *outname*.
2. `PEX_SUFFIX` was not set in *flags*. In this case, if *outname* is not `NULL`, it is used as the output file name. If *outname* is `NULL`, and *tempbase* was not NULL, the output file name is randomly chosen using *tempbase*. Otherwise the output file name is chosen completely at random.

errname is the file name to use for standard error output. If it is `NULL`, standard error is the same as the caller's. Otherwise, standard error is written to the named file.

On an error return, the code sets `*err` to an `errno` value, or to 0 if there is no relevant `errno`.

`const char * pex_run_in_environment (struct pex_obj *obj, int flags, const char *executable, char * const *argv, char * const *env, int env_size, const char *outname, const char *errname, int *err)` [Extension]

Execute one program in a pipeline, permitting the environment for the program to be specified. Behaviour and parameters not listed below are as for `pex_run`.

env is the environment for the child process, specified as an array of character pointers. Each element of the array should point to a string of the form `VAR=VALUE`, with the exception of the last element that must be `NULL`.

`int pexecute (const char *program, char * const *argv, const char *this_pname, const char *temp_base, char **errmsg_fmt, char **errmsg_arg, int flags)` [Extension]

This is the old interface to execute one or more programs. It is still supported for compatibility purposes, but is no longer documented.

`void psignal (int signo, char *message)` [Supplemental]

Print *message* to the standard error, followed by a colon, followed by the description of the signal specified by *signo*, followed by a newline.

`int putenv (const char *string)` [Supplemental]

Uses `setenv` or `unsetenv` to put *string* into the environment or remove it. If *string* is of the form '`name=value`' the string is added; if no '`=`' is present the name is unset/removed.

`int pwait (int pid, int *status, int flags)` [Extension]

Another part of the old execution interface.

`long int random (void)` [Supplement]
`void srandom (unsigned int seed)` [Supplement]
`void* initstate (unsigned int seed, void *arg_state, unsigned long n)` [Supplement]
`void* setstate (void *arg_state)` [Supplement]

Random number functions. `random` returns a random number in the range 0 to `LONG_MAX`. `srandom` initializes the random number generator to some starting point determined by *seed* (else, the values returned by `random` are always the same for each run of the program). `initstate` and `setstate` allow fine-grained control over the state of the random number generator.

`char* reconcat (char *optr, const char *s1, ..., NULL)` [Extension]

Same as `concat`, except that if *optr* is not `NULL` it is freed after the string is created. This is intended to be useful when you're extending an existing string or building up a string in a loop:

```
str = reconcat (str, "pre-", str, NULL);
```

`int rename (const char *old, const char *new)` [Supplemental]

Renames a file from *old* to *new*. If *new* already exists, it is removed.

`char* rindex (const char *s, int c)` [Supplemental]

Returns a pointer to the last occurrence of the character *c* in the string *s*, or `NULL` if not found. The use of `rindex` is deprecated in new programs in favor of `strrchr`.

`int setenv (const char *name, const char *value, int overwrite)` [Supplemental]

`void unsetenv (const char *name)` [Supplemental]

`setenv` adds *name* to the environment with value *value*. If the name was already present in the environment, the new value will be stored only if *overwrite* is nonzero. The companion `unsetenv` function removes *name* from the environment. This implementation is not safe for multithreaded code.

`void setproctitle (const char *fmt, ...)` [Supplemental]

Set the title of a process to *fmt*. va args not supported for now, but defined for compatibility with BSD.

`int signo_max (void)` [Extension]

Returns the maximum signal value for which a corresponding symbolic name or message is available. Note that in the case where we use the `sys_siglist` supplied by the system, it is possible for there to be more symbolic names than messages, or vice versa. In fact, the manual page for `psignal(3b)` explicitly warns that one should check the size of the table (`NSIG`) before indexing it, since new signal codes may be added to the system before they are added to the table. Thus `NSIG` might be smaller than value implied by the largest signo value defined in `<signal.h>`.

We return the maximum value that can be used to obtain a meaningful symbolic name or message.

`int sigsetmask (int set)` [Supplemental]

Sets the signal mask to the one provided in *set* and returns the old mask (which, for libiberty's implementation, will always be the value `1`).

`const char * simple_object_attributes_compare (simple_object_attributes *attrs1, simple_object_attributes *attrs2, int *err)` [Extension]

Compare *attrs1* and *attrs2*. If they could be linked together without error, return `NULL`. Otherwise, return an error message and set `*err` to an errno value or `0` if there is no relevant errno.

`simple_object_attributes * simple_object_fetch_attributes (simple_object_read *simple_object, const char **errmsg, int *err)` [Extension]

Fetch the attributes of *simple_object*. The attributes are internal information such as the format of the object file, or the architecture it was compiled for. This information will persist until `simple_object_attributes_release` is called, even if *simple_object* itself is released.

On error this returns `NULL`, sets `*errmsg` to an error message, and sets `*err` to an errno value or `0` if there is no relevant errno.

`int simple_object_find_section (simple_object_read *simple_object off_t *offset, off_t *length, const char **errmsg, int *err)` [Extension]

Look for the section *name* in *simple_object*. This returns information for the first section with that name.

If found, return 1 and set `*offset` to the offset in the file of the section contents and set `*length` to the length of the section contents. The value in `*offset` will be relative to the offset passed to `simple_object_open_read`.

If the section is not found, and no error occurs, `simple_object_find_section` returns `0` and set `*errmsg` to `NULL`.

If an error occurs, `simple_object_find_section` returns `0`, sets `*errmsg` to an error message, and sets `*err` to an errno value or `0` if there is no relevant errno.

`const char * simple_object_find_sections (simple_object_read *simple_object, int (*pfn) (void *data, const char *name, off_t offset, off_t length), void *data, int *err)` [Extension]

This function calls *pfn* for each section in *simple_object*. It calls *pfn* with the section name, the offset within the file of the section contents, and the length of the section contents. The offset within the file is relative to the offset passed to `simple_object_open_read`. The *data* argument to this function is passed along to *pfn*.

If *pfn* returns `0`, the loop over the sections stops and `simple_object_find_sections` returns. If *pfn* returns some other value, the loop continues.

On success `simple_object_find_sections` returns. On error it returns an error string, and sets `*err` to an errno value or `0` if there is no relevant errno.

`simple_object_read * simple_object_open_read (int descriptor, off_t offset, const char *segment_name, const char **errmsg, int *err)` [Extension]

Opens an object file for reading. Creates and returns an `simple_object_read` pointer which may be passed to other functions to extract data from the object file.

descriptor holds a file descriptor which permits reading.

offset is the offset into the file; this will be `0` in the normal case, but may be a different value when reading an object file in an archive file.

segment_name is only used with the Mach-O file format used on Darwin aka Mac OS X. It is required on that platform, and means to only look at sections within the segment with that name. The parameter is ignored on other systems.

If an error occurs, this functions returns `NULL` and sets `*errmsg` to an error string and sets `*err` to an errno value or `0` if there is no relevant errno.

`void simple_object_release_attributes (simple_object_attributes *attrs)` [Extension]

Release all resources associated with *attrs*.

`void simple_object_release_read (simple_object_read *simple_object)` [Extension]

Release all resources associated with *simple_object*. This does not close the file descriptor.

`void simple_object_release_write (simple_object_write *simple_object)` [Extension]

Release all resources associated with *simple_object*.

`simple_object_write * simple_object_start_write (simple_object_attributes attrs, const char *segment_name, const char **errmsg, int *err)` [Extension]

Start creating a new object file using the object file format described in *attrs*. You must fetch attribute information from an existing object file before you can create a new one. There is currently no support for creating an object file de novo.

segment_name is only used with Mach-O as found on Darwin aka Mac OS X. The parameter is required on that target. It means that all sections are created within the named segment. It is ignored for other object file formats.

On error `simple_object_start_write` returns `NULL`, sets `*ERRMSG` to an error message, and sets `*err` to an errno value or `0` if there is no relevant errno.

`const char * simple_object_write_add_data (simple_object_write *simple_object, simple_object_write_section *section, const void *buffer, size_t size, int copy, int *err)` [Extension]

Add data *buffer*/*size* to *section* in *simple_object*. If *copy* is non-zero, the data will be copied into memory if necessary. If *copy* is zero, *buffer* must persist until `simple_object_write_to_file` is called. is released.

On success this returns `NULL`. On error this returns an error message, and sets `*err` to an errno value or 0 if there is no relevant erro.

`simple_object_write_section * simple_object_write_create_section (simple_object_write *simple_object, const char *name, unsigned int align, const char **errmsg, int *err)` [Extension]

Add a section to *simple_object*. *name* is the name of the new section. *align* is the required alignment expressed as the number of required low-order 0 bits (e.g., 2 for alignment to a 32-bit boundary).

The section is created as containing data, readable, not writable, not executable, not loaded at runtime. The section is not written to the file until `simple_object_write_to_file` is called.

On error this returns `NULL`, sets `*errmsg` to an error message, and sets `*err` to an errno value or `0` if there is no relevant errno.

`const char * simple_object_write_to_file (simple_object_write *simple_object, int descriptor, int *err)` [Extension]

Write the complete object file to *descriptor*, an open file descriptor. This writes out all the data accumulated by calls to `simple_object_write_create_section` and *simple_object_write_add_data*.

This returns `NULL` on success. On error this returns an error message and sets `*err` to an errno value or `0` if there is no relevant errno.

`int snprintf (char *buf, size_t n, const char *format, ...)` [Supplemental]

This function is similar to `sprintf`, but it will write to *buf* at most `n-1` bytes of text, followed by a terminating null byte, for a total of *n* bytes. On error the return value is -1, otherwise it returns the number of bytes, not including the terminating null byte, that would have been written had *n* been sufficiently large, regardless of the actual value of *n*. Note some pre-C99 system libraries do not implement this correctly so users cannot generally rely on the return value if the system version of this function is used.

`char* spaces (int count)` [Extension]

Returns a pointer to a memory region filled with the specified number of spaces and null terminated. The returned pointer is valid until at least the next call.

`splay_tree splay_tree_new_with_typed_alloc (splay_tree_compare_fn compare_fn, splay_tree_delete_key_fn delete_key_fn, splay_tree_delete_value_fn delete_value_fn, splay_tree_allocate_fn tree_allocate_fn, splay_tree_allocate_fn node_allocate_fn, splay_tree_deallocate_fn deallocate_fn, void * allocate_data)` [Supplemental]

This function creates a splay tree that uses two different allocators *tree_allocate_fn* and *node_allocate_fn* to use for allocating the tree itself and its nodes respectively. This is useful when variables of different types need to be allocated with different allocators.

The splay tree will use *compare_fn* to compare nodes, *delete_key_fn* to deallocate keys, and *delete_value_fn* to deallocate values. Keys and values will be deallocated when the tree is deleted using splay_tree_delete or when a node is removed using splay_tree_remove. splay_tree_insert will release the previously inserted key and value using *delete_key_fn* and *delete_value_fn* if the inserted key is already found in the tree.

`void stack_limit_increase (unsigned long pref)` [Extension]

Attempt to increase stack size limit to *pref* bytes if possible.

`char* stpcpy (char *dst, const char *src)` [Supplemental]

Copies the string *src* into *dst*. Returns a pointer to *dst* + strlen(*src*).

`char* stpncpy (char *dst, const char *src, size_t len)` [Supplemental]

Copies the string *src* into *dst*, copying exactly *len* and padding with zeros if necessary. If *len* < strlen(*src*) then return *dst* + *len*, otherwise returns *dst* + strlen(*src*).

`int strcasecmp (const char *s1, const char *s2)` [Supplemental]
A case-insensitive `strcmp`.

`char* strchr (const char *s, int c)` [Supplemental]
Returns a pointer to the first occurrence of the character *c* in the string *s*, or `NULL` if not found. If *c* is itself the null character, the results are undefined.

`char* strdup (const char *s)` [Supplemental]
Returns a pointer to a copy of *s* in memory obtained from `malloc`, or `NULL` if insufficient memory was available.

`const char* strerrno (int errnum)` [Replacement]
Given an error number returned from a system call (typically returned in `errno`), returns a pointer to a string containing the symbolic name of that error number, as found in `<errno.h>`.

If the supplied error number is within the valid range of indices for symbolic names, but no name is available for the particular error number, then returns the string '`Error num`', where *num* is the error number.

If the supplied error number is not within the range of valid indices, then returns `NULL`.

The contents of the location pointed to are only guaranteed to be valid until the next call to `strerrno`.

`char* strerror (int errnoval)` [Supplemental]
Maps an `errno` number to an error message string, the contents of which are implementation defined. On systems which have the external variables `sys_nerr` and `sys_errlist`, these strings will be the same as the ones used by `perror`.

If the supplied error number is within the valid range of indices for the `sys_errlist`, but no message is available for the particular error number, then returns the string '`Error num`', where *num* is the error number.

If the supplied error number is not a valid index into `sys_errlist`, returns `NULL`.

The returned string is only guaranteed to be valid only until the next call to `strerror`.

`int strncasecmp (const char *s1, const char *s2)` [Supplemental]
A case-insensitive `strncmp`.

`int strncmp (const char *s1, const char *s2, size_t n)` [Supplemental]
Compares the first *n* bytes of two strings, returning a value as `strcmp`.

`char* strndup (const char *s, size_t n)` [Extension]
Returns a pointer to a copy of *s* with at most *n* characters in memory obtained from `malloc`, or `NULL` if insufficient memory was available. The result is always NUL terminated.

`size_t strnlen (const char *s, size_t maxlen)` [Supplemental]
Returns the length of *s*, as with `strlen`, but never looks past the first *maxlen* characters in the string. If there is no '\0' character in the first *maxlen* characters, returns *maxlen*.

`char* strrchr (const char *s, int c)` [Supplemental]

Returns a pointer to the last occurrence of the character *c* in the string *s*, or `NULL` if not found. If *c* is itself the null character, the results are undefined.

`const char * strsignal (int signo)` [Supplemental]

Maps an signal number to an signal message string, the contents of which are implementation defined. On systems which have the external variable `sys_siglist`, these strings will be the same as the ones used by `psignal()`.

If the supplied signal number is within the valid range of indices for the `sys_siglist`, but no message is available for the particular signal number, then returns the string '`Signal num`', where *num* is the signal number.

If the supplied signal number is not a valid index into `sys_siglist`, returns `NULL`.

The returned string is only guaranteed to be valid only until the next call to `strsignal`.

`const char* strsigno (int signo)` [Extension]

Given an signal number, returns a pointer to a string containing the symbolic name of that signal number, as found in `<signal.h>`.

If the supplied signal number is within the valid range of indices for symbolic names, but no name is available for the particular signal number, then returns the string '`Signal num`', where *num* is the signal number.

If the supplied signal number is not within the range of valid indices, then returns `NULL`.

The contents of the location pointed to are only guaranteed to be valid until the next call to `strsigno`.

`char* strstr (const char *string, const char *sub)` [Supplemental]

This function searches for the substring *sub* in the string *string*, not including the terminating null characters. A pointer to the first occurrence of *sub* is returned, or `NULL` if the substring is absent. If *sub* points to a string with zero length, the function returns *string*.

`double strtod (const char *string, char **endptr)` [Supplemental]

This ISO C function converts the initial portion of *string* to a `double`. If *endptr* is not `NULL`, a pointer to the character after the last character used in the conversion is stored in the location referenced by *endptr*. If no conversion is performed, zero is returned and the value of *string* is stored in the location referenced by *endptr*.

`int strtoerrno (const char *name)` [Extension]

Given the symbolic name of a error number (e.g., `EACCES`), map it to an errno value. If no translation is found, returns 0.

`long int strtol (const char *string, char **endptr, int base)` [Supplemental]

`unsigned long int strtoul (const char *string, char **endptr, int base)` [Supplemental]

The `strtol` function converts the string in *string* to a long integer value according to the given *base*, which must be between 2 and 36 inclusive, or be the special value

0. If *base* is 0, `strtol` will look for the prefixes `0` and `0x` to indicate bases 8 and 16, respectively, else default to base 10. When the base is 16 (either explicitly or implicitly), a prefix of `0x` is allowed. The handling of *endptr* is as that of `strtod` above. The `strtoul` function is the same, except that the converted value is unsigned.

`long long int strtoll (const char *string, char **endptr, int base)` [Supplemental]

`unsigned long long int strtoull ( const char *string, char **endptr, int base)` [Supplemental]

The `strtoll` function converts the string in *string* to a long long integer value according to the given *base*, which must be between 2 and 36 inclusive, or be the special value 0. If *base* is 0, `strtoll` will look for the prefixes `0` and `0x` to indicate bases 8 and 16, respectively, else default to base 10. When the base is 16 (either explicitly or implicitly), a prefix of `0x` is allowed. The handling of *endptr* is as that of `strtod` above. The `strtoull` function is the same, except that the converted value is unsigned.

`int strtosigno (const char *name)` [Extension]

Given the symbolic name of a signal, map it to a signal number. If no translation is found, returns 0.

`int strverscmp (const char *s1, const char *s2)` [Function]

The `strverscmp` function compares the string *s1* against *s2*, considering them as holding indices/version numbers. Return value follows the same conventions as found in the `strverscmp` function. In fact, if *s1* and *s2* contain no digits, `strverscmp` behaves like `strcmp`.

Basically, we compare strings normally (character by character), until we find a digit in each string - then we enter a special comparison mode, where each sequence of digits is taken as a whole. If we reach the end of these two parts without noticing a difference, we return to the standard comparison mode. There are two types of numeric parts: "integral" and "fractional" (those begin with a '0'). The types of the numeric parts affect the way we sort them:

- integral/integral: we compare values as you would expect.
- fractional/integral: the fractional part is less than the integral one. Again, no surprise.
- fractional/fractional: the things become a bit more complex. If the common prefix contains only leading zeroes, the longest part is less than the other one; else the comparison behaves normally.

```
strverscmp ("no digit", "no digit")
    ⇒ 0    // same behavior as strcmp.
strverscmp ("item#99", "item#100")
    ⇒ <0   // same prefix, but 99 < 100.
strverscmp ("alpha1", "alpha001")
    ⇒ >0   // fractional part inferior to integral one.
strverscmp ("part1_f012", "part1_f01")
    ⇒ >0   // two fractional parts.
strverscmp ("foo.009", "foo.0")
    ⇒ <0   // idem, but with leading zeroes only.
```

This function is especially useful when dealing with filename sorting, because filenames frequently hold indices/version numbers.

`void timeval_add (struct timeval *a, struct timeval *b, struct timeval *result)` [Extension]

Adds *a* to *b* and stores the result in *result*.

`void timeval_sub (struct timeval *a, struct timeval *b, struct timeval *result)` [Extension]

Subtracts *b* from *a* and stores the result in *result*.

`char* tmpnam (char *s)` [Supplemental]

This function attempts to create a name for a temporary file, which will be a valid file name yet not exist when `tmpnam` checks for it. *s* must point to a buffer of at least `L_tmpnam` bytes, or be `NULL`. Use of this function creates a security risk, and it must not be used in new projects. Use `mkstemp` instead.

`int unlink_if_ordinary (const char*)` [Supplemental]

Unlinks the named file, unless it is special (e.g. a device file). Returns 0 when the file was unlinked, a negative value (and errno set) when there was an error deleting the file, and a positive value if no attempt was made to unlink the file because it is special.

`void unlock_std_streams (void)` [Extension]

If the OS supports it, ensure that the standard I/O streams, `stdin`, `stdout` and `stderr` are setup to avoid any multi-threaded locking. Otherwise do nothing.

`void unlock_stream (FILE * stream)` [Extension]

If the OS supports it, ensure that the supplied stream is setup to avoid any multi-threaded locking. Otherwise leave the `FILE` pointer unchanged. If the *stream* is `NULL` do nothing.

`int vasprintf (char **resptr, const char *format, va_list args)` [Extension]

Like `vsprintf`, but instead of passing a pointer to a buffer, you pass a pointer to a pointer. This function will compute the size of the buffer needed, allocate memory with `malloc`, and store a pointer to the allocated memory in `*resptr`. The value returned is the same as `vsprintf` would return. If memory could not be allocated, minus one is returned and `NULL` is stored in `*resptr`.

`int vfork (void)` [Supplemental]

Emulates `vfork` by calling `fork` and returning its value.

`int vprintf (const char *format, va_list ap)` [Supplemental]
`int vfprintf (FILE *stream, const char *format, va_list ap)` [Supplemental]
`int vsprintf (char *str, const char *format, va_list ap)` [Supplemental]

These functions are the same as `printf`, `fprintf`, and `sprintf`, respectively, except that they are called with a `va_list` instead of a variable number of arguments. Note that they do not call `va_end`; this is the application's responsibility. In `libiberty` they are implemented in terms of the nonstandard but common function `_doprnt`.

`int vsnprintf (char *buf, size_t n, const char *format, va_list ap)` [Supplemental]

This function is similar to `vsprintf`, but it will write to *buf* at most `n-1` bytes of text, followed by a terminating null byte, for a total of *n* bytes. On error the return value is -1, otherwise it returns the number of characters that would have been printed had *n* been sufficiently large, regardless of the actual value of *n*. Note some pre-C99 system libraries do not implement this correctly so users cannot generally rely on the return value if the system version of this function is used.

`int waitpid (int pid, int *status, int)` [Supplemental]

This is a wrapper around the `wait` function. Any "special" values of *pid* depend on your implementation of `wait`, as does the return value. The third argument is unused in `libiberty`.

`int writeargv (char * const *argv, FILE *file)` [Extension]

Write each member of ARGV, handling all necessary quoting, to the file named by FILE, separated by whitespace. Return 0 on success, non-zero if an error occurred while writing to FILE.

`char* xasprintf (const char *format, ...)` [Replacement]

Print to allocated string without fail. If `xasprintf` fails, this will print a message to `stderr` (using the name set by `xmalloc_set_program_name`, if any) and then call `xexit`.

`int xatexit (void (*fn) (void))` [Function]

Behaves as the standard `atexit` function, but with no limit on the number of registered functions. Returns 0 on success, or -1 on failure. If you use `xatexit` to register functions, you must use `xexit` to terminate your program.

`void* xcalloc (size_t nelem, size_t elsize)` [Replacement]

Allocate memory without fail, and set it to zero. This routine functions like `calloc`, but will behave the same as `xmalloc` if memory cannot be found.

`void xexit (int code)` [Replacement]

Terminates the program. If any functions have been registered with the `xatexit` replacement function, they will be called first. Termination is handled via the system's normal `exit` call.

`void* xmalloc (size_t)` [Replacement]

Allocate memory without fail. If `malloc` fails, this will print a message to `stderr` (using the name set by `xmalloc_set_program_name`, if any) and then call `xexit`. Note that it is therefore safe for a program to contain `#define malloc xmalloc` in its source.

`void xmalloc_failed (size_t)` [Replacement]

This function is not meant to be called by client code, and is listed here for completeness only. If any of the allocation routines fail, this function will be called to print an error message and terminate execution.

`void xmalloc_set_program_name (const char *name)` [Replacement]

You can use this to set the name of the program used by `xmalloc_failed` when printing a failure message.

`void* xmemdup (void *input, size_t copy_size, size_t alloc_size)` [Replacement]

Duplicates a region of memory without fail. First, *alloc_size* bytes are allocated, then *copy_size* bytes from *input* are copied into it, and the new memory is returned. If fewer bytes are copied than were allocated, the remaining memory is zeroed.

`void* xrealloc (void *ptr, size_t size)` [Replacement]

Reallocate memory without fail. This routine functions like `realloc`, but will behave the same as `xmalloc` if memory cannot be found.

`char* xstrdup (const char *s)` [Replacement]

Duplicates a character string without fail, using `xmalloc` to obtain memory.

`char* xstrerror (int errnum)` [Replacement]

Behaves exactly like the standard `strerror` function, but will never return a `NULL` pointer.

`char* xstrndup (const char *s, size_t n)` [Replacement]

Returns a pointer to a copy of *s* with at most *n* characters without fail, using `xmalloc` to obtain memory. The result is always NUL terminated.

`char* xvasprintf (const char *format, va_list args)` [Replacement]

Print to allocated string without fail. If `xvasprintf` fails, this will print a message to `stderr` (using the name set by `xmalloc_set_program_name`, if any) and then call `xexit`.

Appendix A Licenses

A.1 GNU LESSER GENERAL PUBLIC LICENSE

Version 2.1, February 1999

[This is the first released version of the Lesser GPL. It also counts
as the successor of the GNU Library Public License, version 2, hence the
version number 2.1.]

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A.2 BSD

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